

```

procedure INIT-PDA
{Invoked when the router comes up. }
begin
    Initialize all tables;
    call PDA;
end INIT-PDA

procedure PDA
{Executed at each router i. Invoked when an event occurs}
begin
    (1) call NTU;
    (2) call MTU; /* Updates  $T^i$  */
    (3) if (there are changes to  $T^i$ ) then
        Compose an LSU message consisting of topology
        differences using add, delete
        and change link entries;
        end if
    (4) Within a finite amount time, send the
        LSU message to all neighbors;
end PDA

```

FIG. 1

```

procedure NTU
begin
    (1) if (LSU message is received from a neighbor  $k$ ) then
        (1a) Update neighbor table  $T_k^i$ . That is, add links,
              delete links or change links according to the
              specification of each entry in the LSU;
        (1b) Run Dijkstra's shortest path algorithm
              on the resulting topology  $T_k^i$ ; /*This results in
              finding minimum distances from  $k$  to all other
              nodes in  $T_k^i$ . Note  $T_k^i$  is a tree*/
        (1c) Update  $D_{jk}^i$  with new distances in  $T_k^i$ ;
    end if
    (2) if (adjacent link  $(i, k)$  is up) then
        Update  $l_k^i$  and send an LSU message to the
        neighbor  $k$  with link information of all links in
        its main topology table  $T^i$ ;
    endif
    (3) if (cost of an adjacent link  $(i, k)$  changed)then
        Update  $l_k^i$ ;
    endif
    (4) if (adjacent link  $(i, k)$  failed)then
        Update  $l_k^i$  and clear the table  $T_k^i$ ;
    endif
end NTU

```

FIG. 2

```

procedure MTU at router  $i$ 
begin
    (1)  $oldT^i \leftarrow T^i$ ; /* Save copy */
    (2) if (node  $j$  occurs in at least one of  $T_k^i$ ) then
        add  $j$  to the main topology table  $T^i$ 
        end if
    (3) for each node  $j$  in  $T^i$  do
         $MIN \leftarrow \min \{D_{jk}^i + l_k^i | k \in N^i\}$ ;
        let  $p$  be such that  $MIN = (D_{jp}^i + l_p^i)$ ;
        /* Neighbor  $p$  is the preferred neighbor for
        destination  $j$ . Ties are broken in favor of
        lower address neighbor */
        done
    (4) for each  $j$  in  $T^i$  and its preferred neighbor  $p$  do
        Copy all links  $(j,n)$  from  $T_p^i$  to  $T^i$ ,
        /* i.e., copy all links in  $T_p^i$  for which
         $j$  is the head node */
        done
    (5) Update  $T^i$  with information of each  $l_k^i$ ;
    (6) Run Dijkstra's shortest path algorithm on  $T^i$ 
        and remove those links in  $T^i$  that are not
        part of the shortest path tree;
    (7) Update  $D_j^i$  with new distances in  $T^i$ ;
    (8) Compare  $oldT^i$  with  $T^i$  and note all differences;
end MTU

```

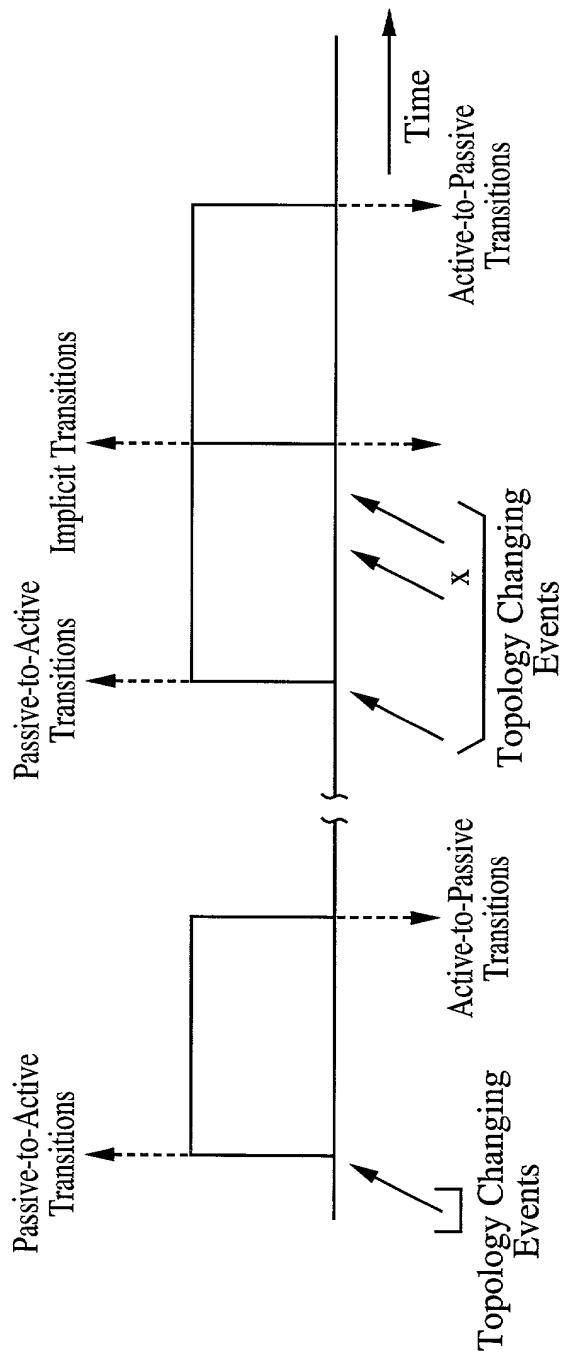
FIG. 3

```

procedure MPDA at router  $i$ 
{invoked when an event occurs}
begin
    (1) call NTU;
    (2) if (node is in PASSIVE state) then
        (2a) call MTU; /* update  $T^i$  and  $D_j^i$  */
        (2b)  $F D_j^i \leftarrow \min\{F D_{j'}^i, D_j^i\}$ ;
        endif
    (3) if (node is in ACTIVE state and the
          last ACK is received) then
        (3a)  $temp_j^i \leftarrow D_j^i$ ; set node to PASSIVE state;
        (3b) call MTU to update  $T^i$ ;
        (3c)  $F D_j^i \leftarrow \min\{temp_{j'}^i, D_j^i\}$ 
        endif
    (4)  $S_j^i \leftarrow \{k | D_j^i < F D_j^i\}$ ;
    (5) if (changes occur in  $T^i$ )then
        Set node to ACTIVE state;
        endif
        if (no changes occur in  $T^i$  and the event is
            the last ACK) then
            Set node to PASSIVE state;
            endif
    (6) if (there are changes to  $T^i$ ) then
        Compose anew LSU with the topology
        changes expressed as add link,
        delete link and change link;
        end if
    (7) if (input event received is an LSU message)then
        Add the ACK entry to newly composed LSU;
        endif
    (8) Send the new LSU message.
end MPDA

```

FIG. 4

**FIG. 5**

www.ijerph.org | ISSN: 1660-4601 | DOI: 10.3390/ijerph16010001 | Cite this article: https://doi.org/10.3390/ijerph16010001

```

Procedure IH
begin
    (1)  $\forall k \notin S_j^i \phi_{jk}^i \leftarrow 0$ ;
    (2) if ( $|S_j^i| = 1$ ) then
         $\forall k \in S_j^i \phi_{jk}^i \leftarrow 1$ ;
    endif
    (3) if ( $|S_j^i| > 1$ ) then

```

$$\phi_{jk}^i \leftarrow \frac{D_{jk}^i + l_k^i}{\sum_{m \in S_j^i} (D_j^i + l_m^i)}, \quad \forall k \in S_j^i;$$

```
endif
end IH
```

FIG. 6

Procedure AH

begin

(1) $D_{min}^{ij} \leftarrow \min \{D_{jk}^i + l_k^i | k \in S_j\};$

(2) let $D_{min}^{ij} = (D_{jk_0}^i + l_{k_0}^i)$
 // that is, k_0 be the neighbor
 that offers the minimum)

(3) foreach $k \in S_j^i$ do

$a_{jk}^i \leftarrow D_{jk}^i + l_k^i - D_{min}^{ij};$

done

(4) $\Delta \leftarrow \frac{1}{2} \min \{ \frac{\phi_{jk}^i}{a_{jk}^i} | k \in S_j^i \wedge a_{jk}^i \neq 0 \};$

(5) foreach $k \neq k_0 \wedge k \in S_j^i$ do

$\phi_{jk}^i \leftarrow \phi_{jk}^i - \Delta \times a_{jk}^i;$

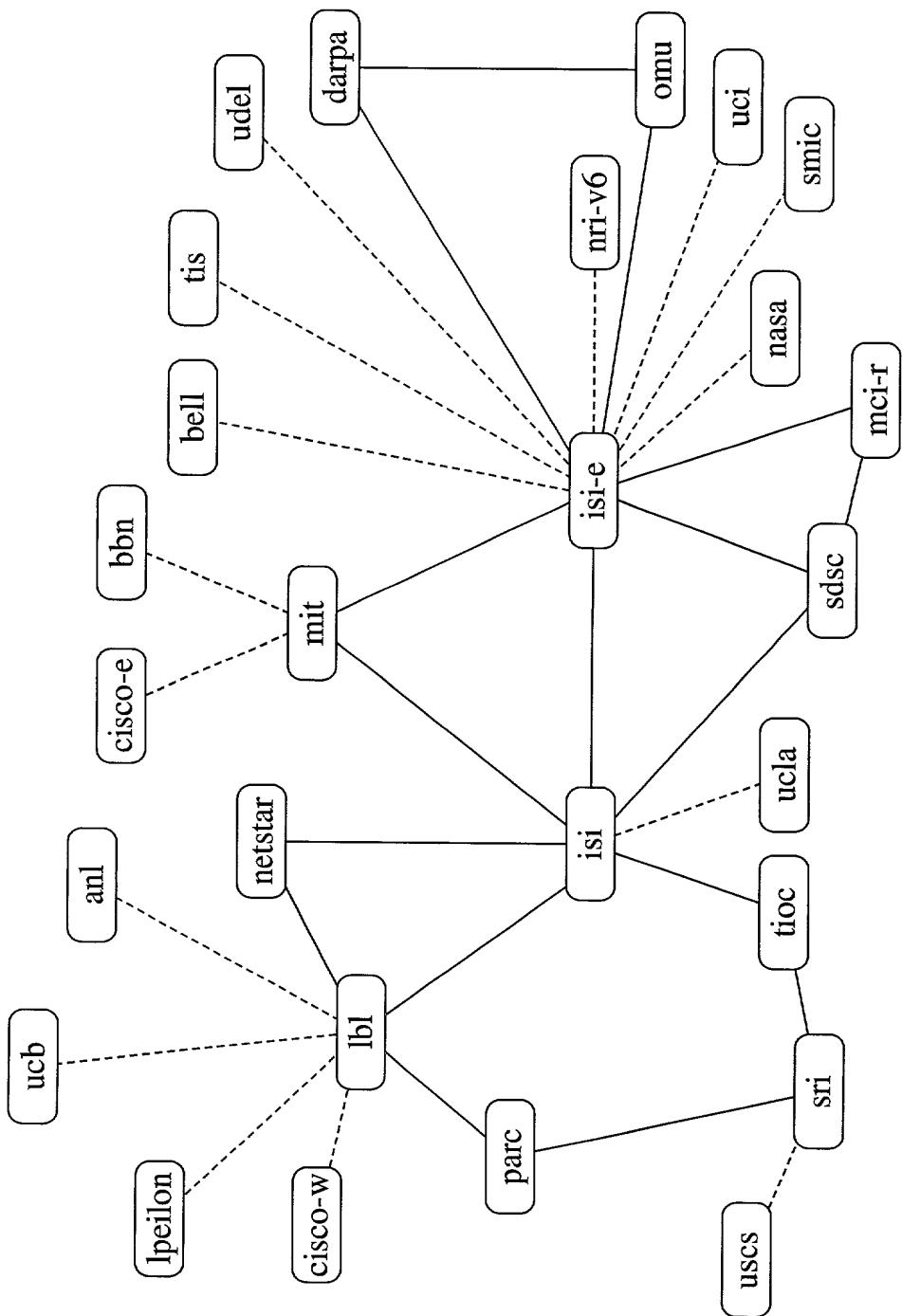
done

(6) foreach $k = k_0$ do

$\phi_{jk}^i \leftarrow \phi_{jk}^i + \sum_{q \in S_k} \Delta \times a_{jq}^i;$
 done

end AH

FIG. 7



8
FIG.

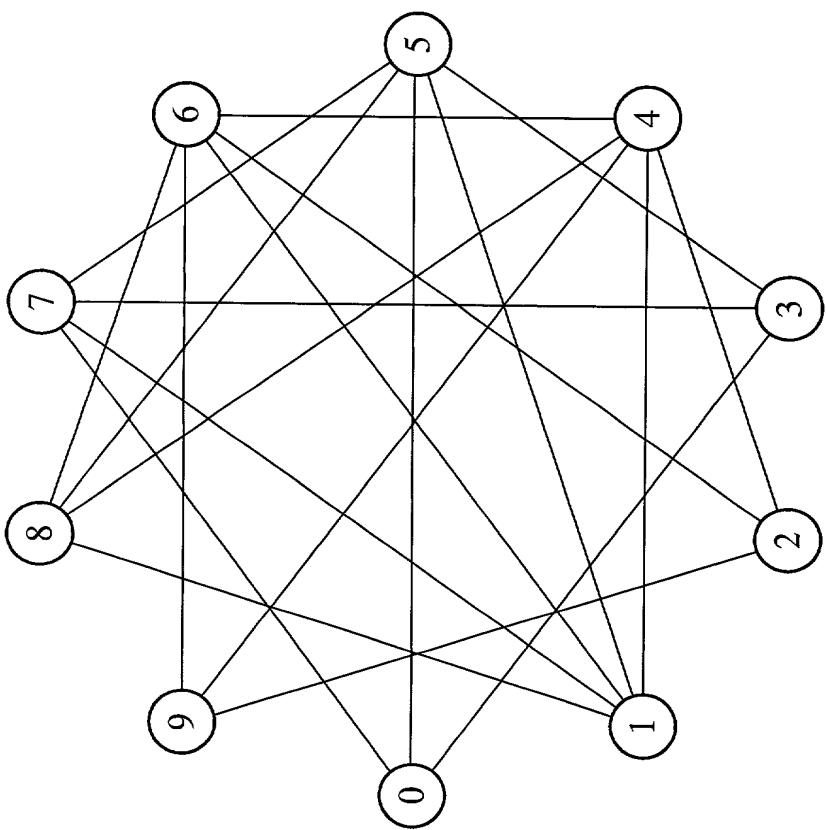
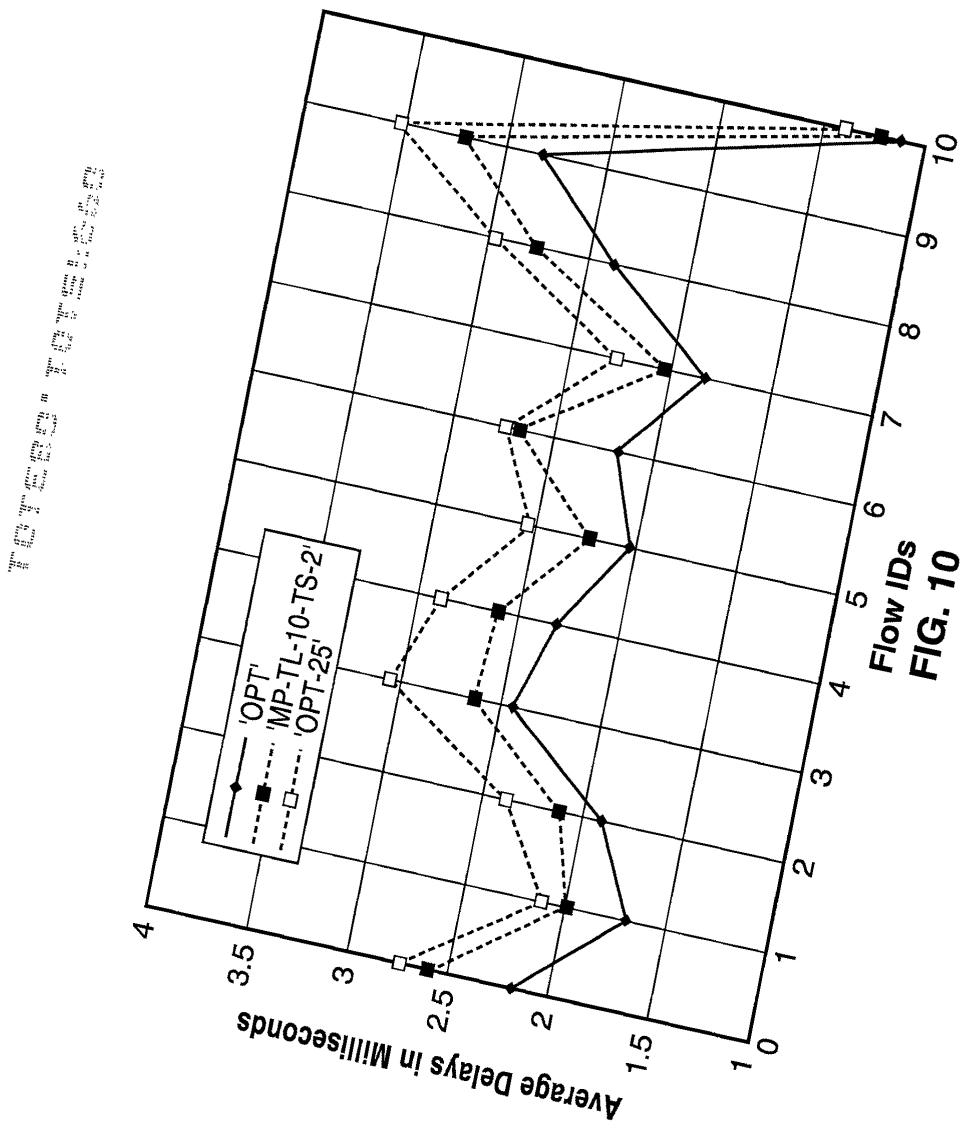


FIG. 9



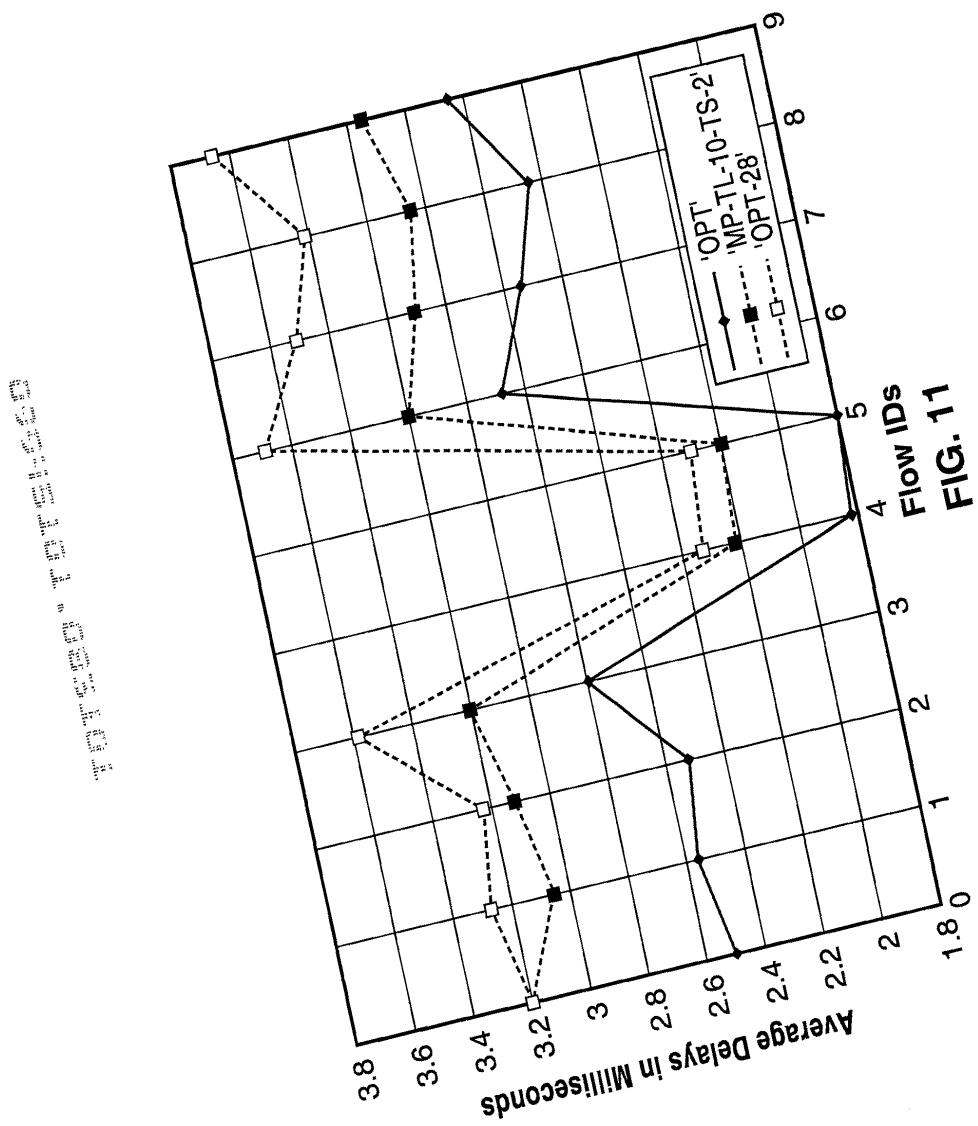


FIG. 12

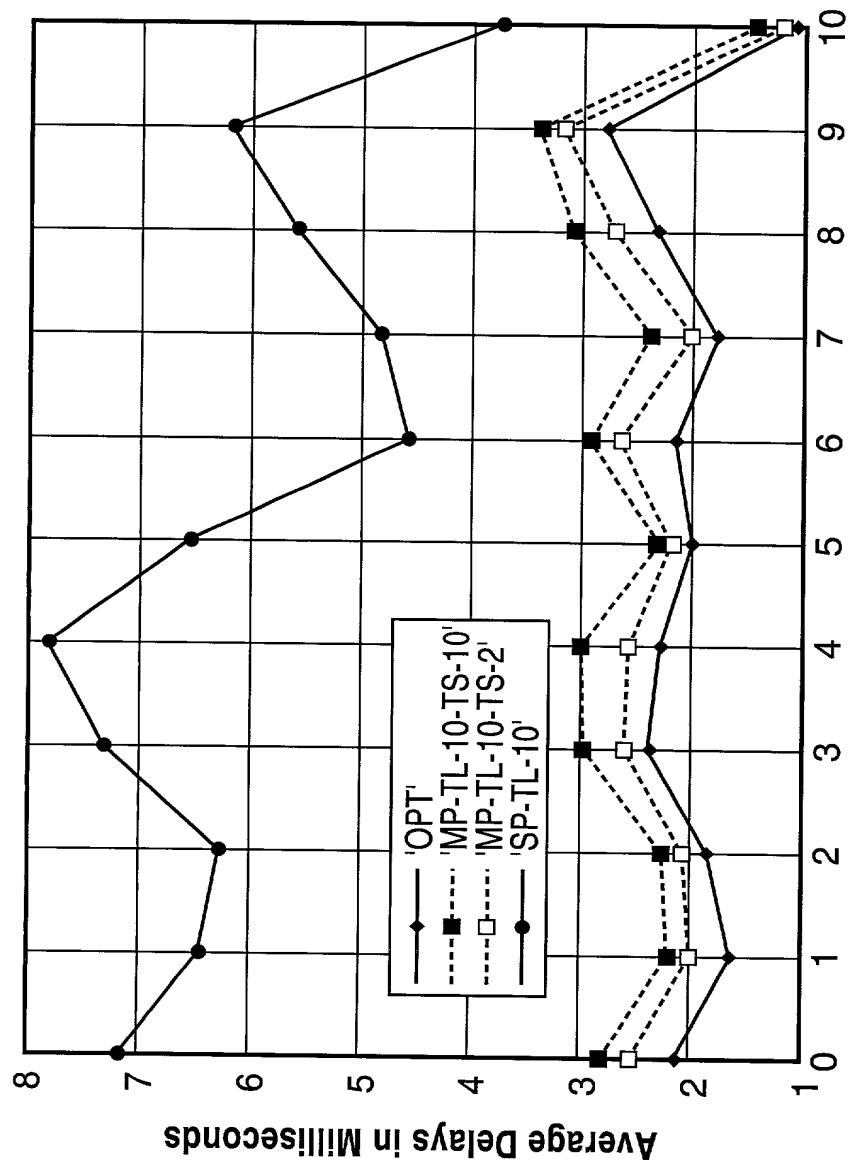
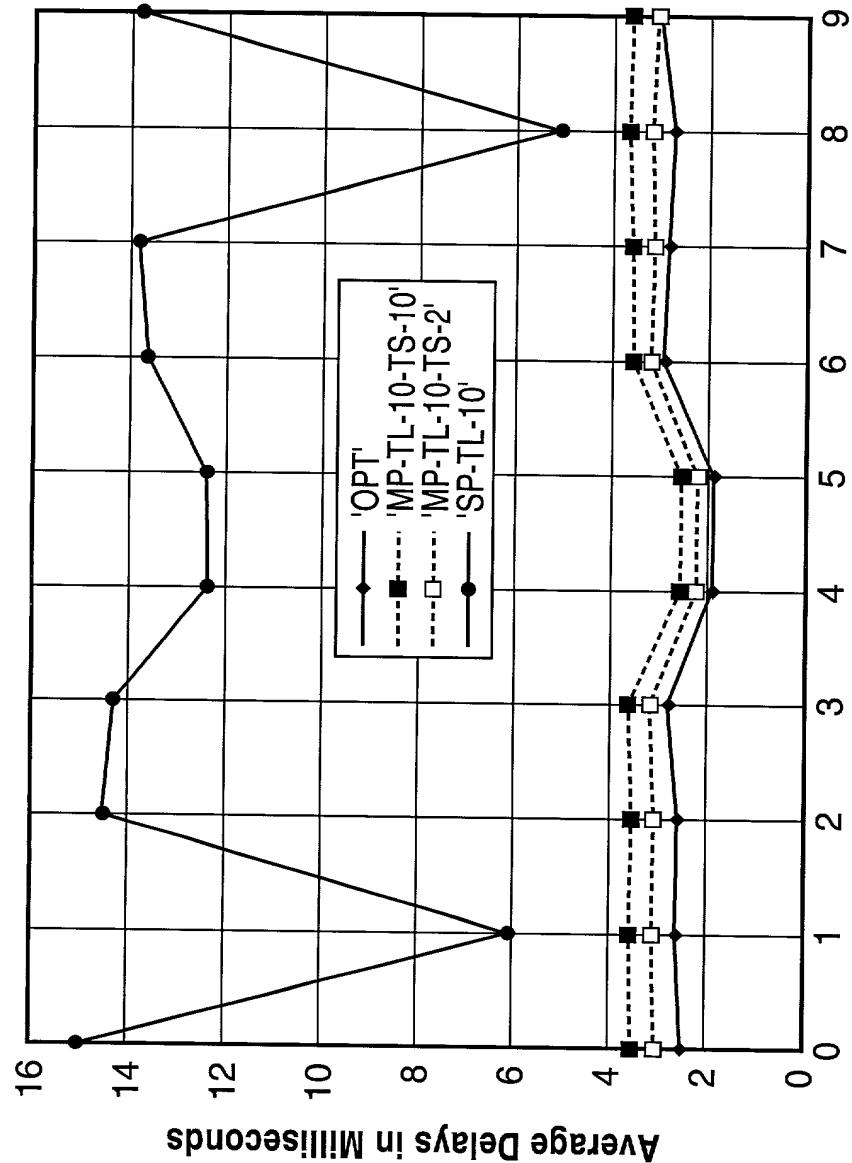


FIG. 13

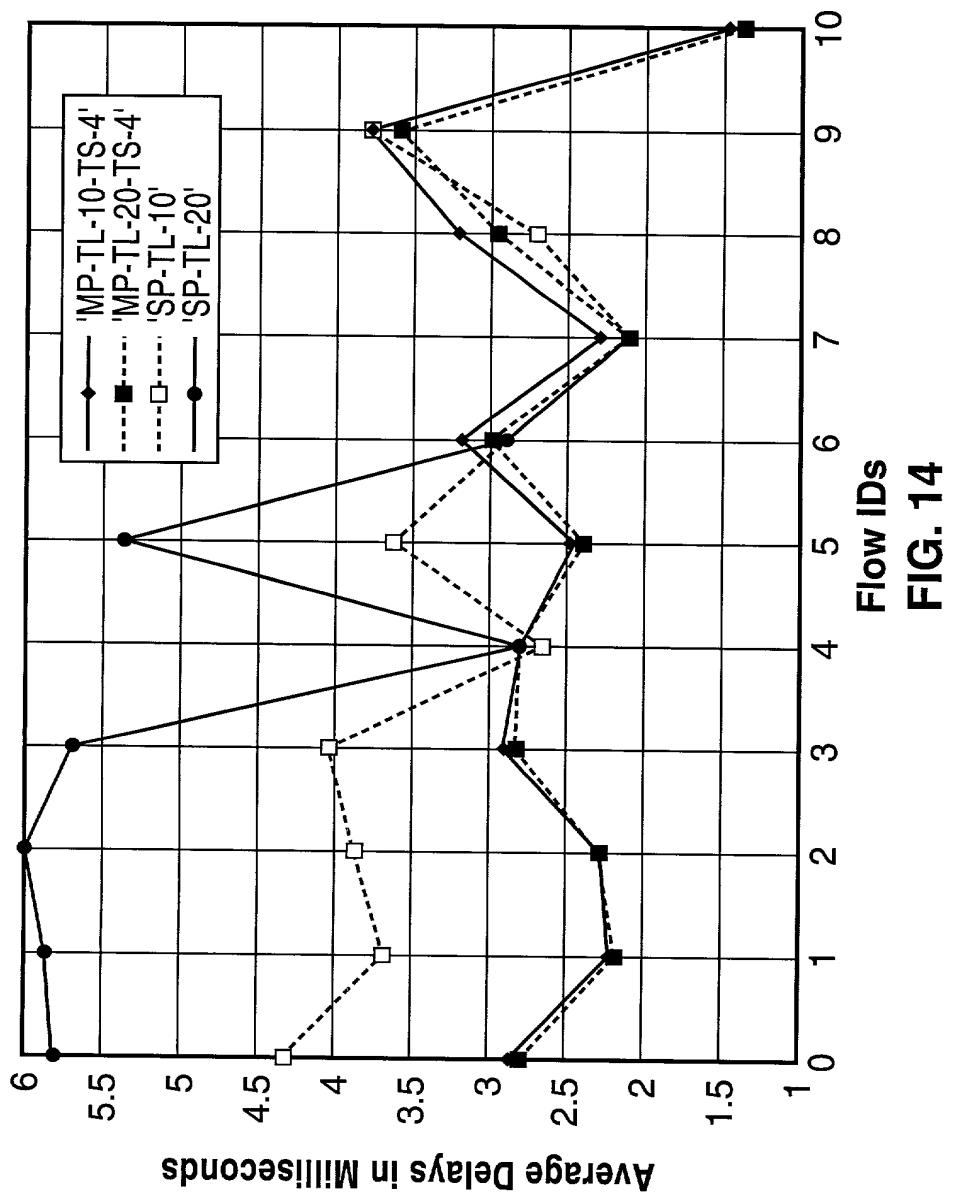
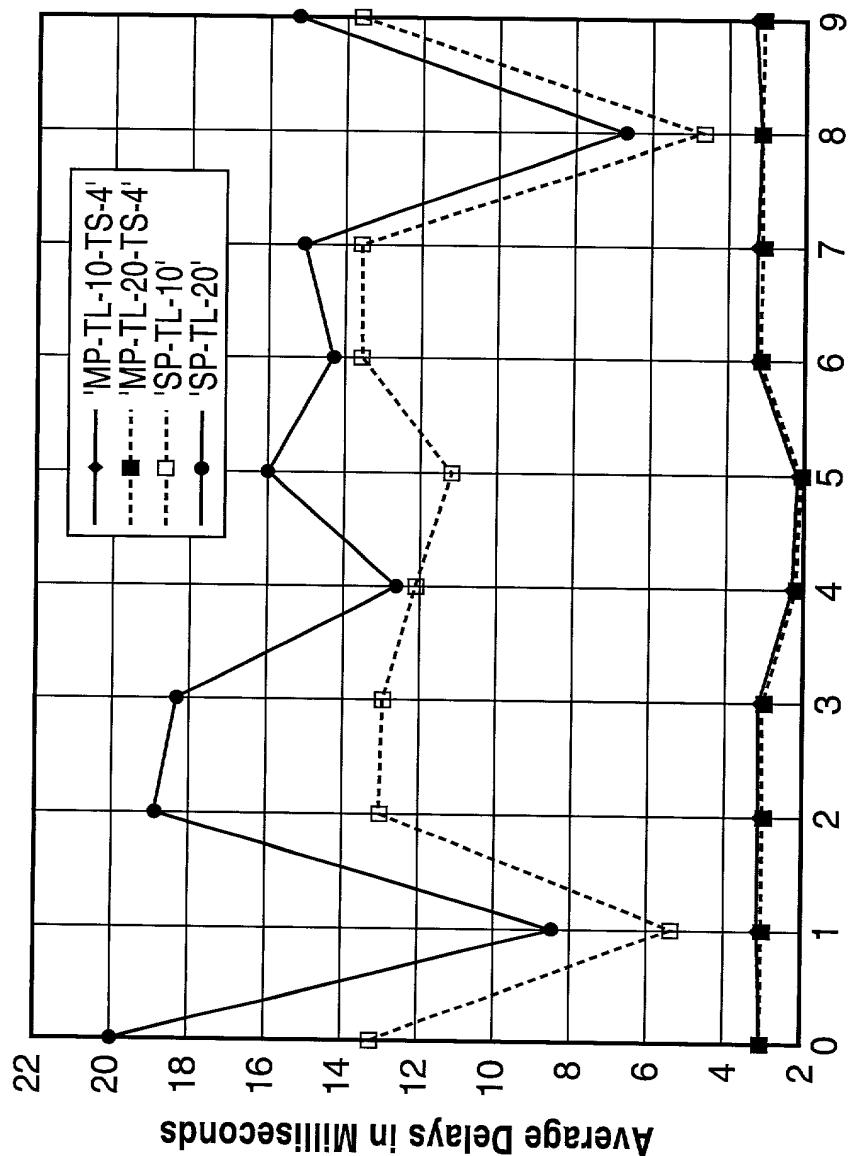


FIG. 15

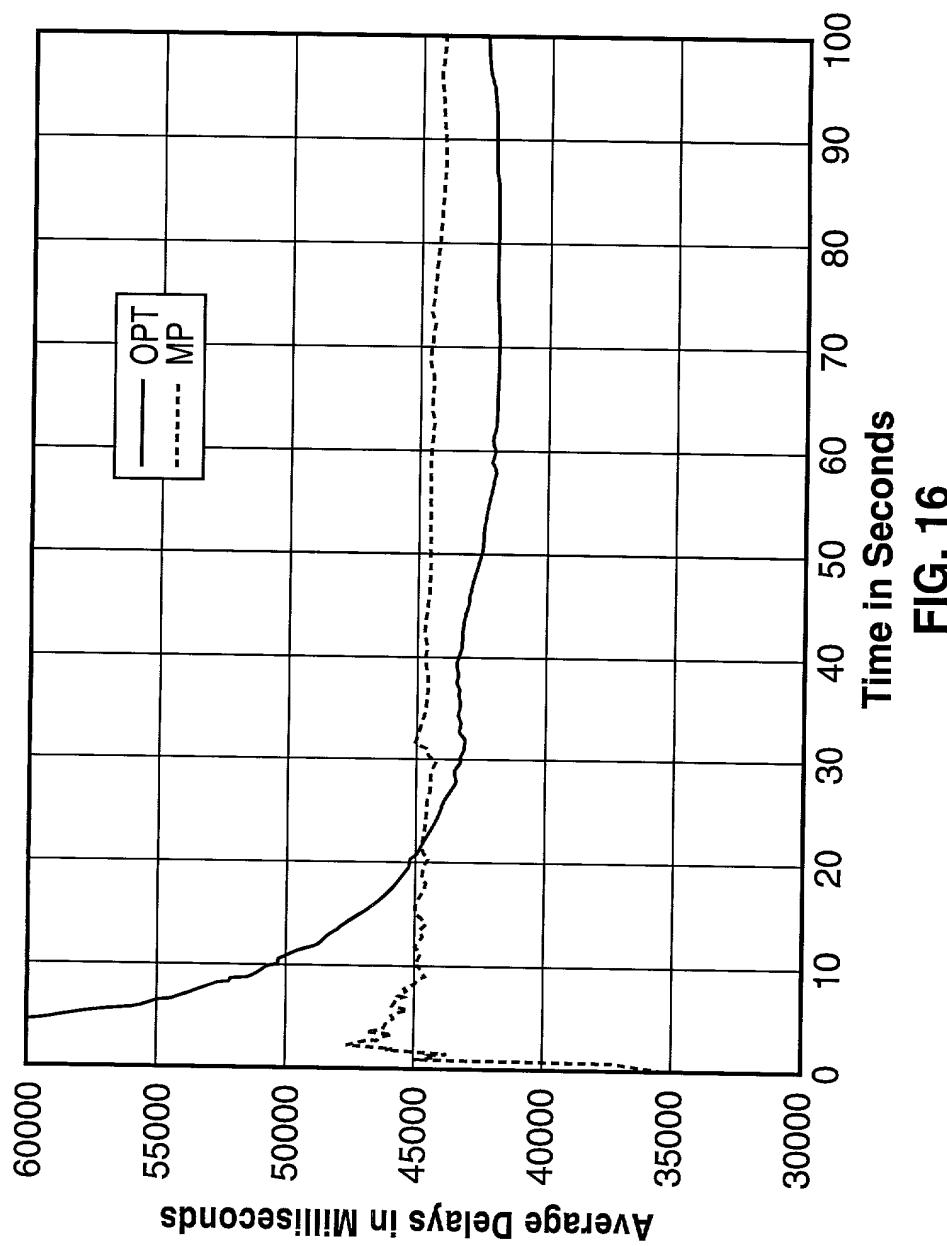
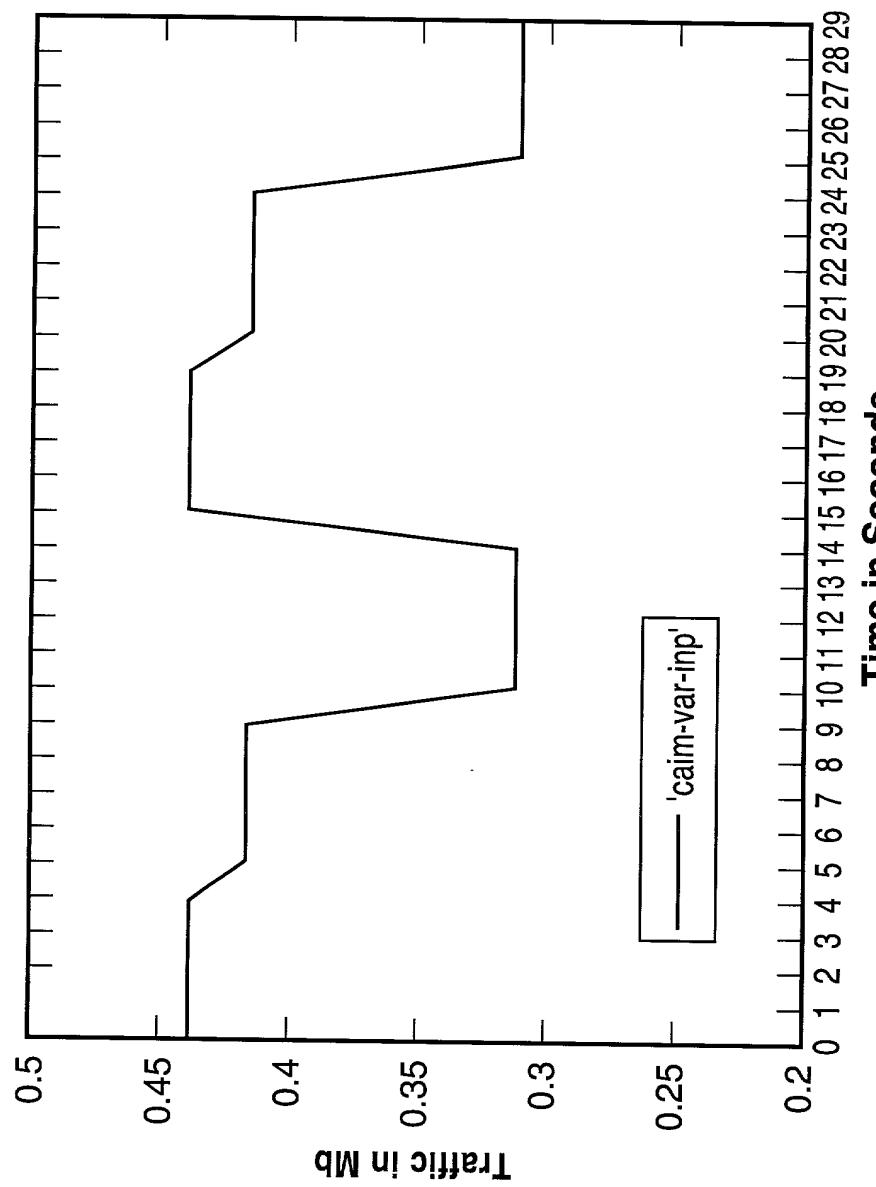


FIG. 17

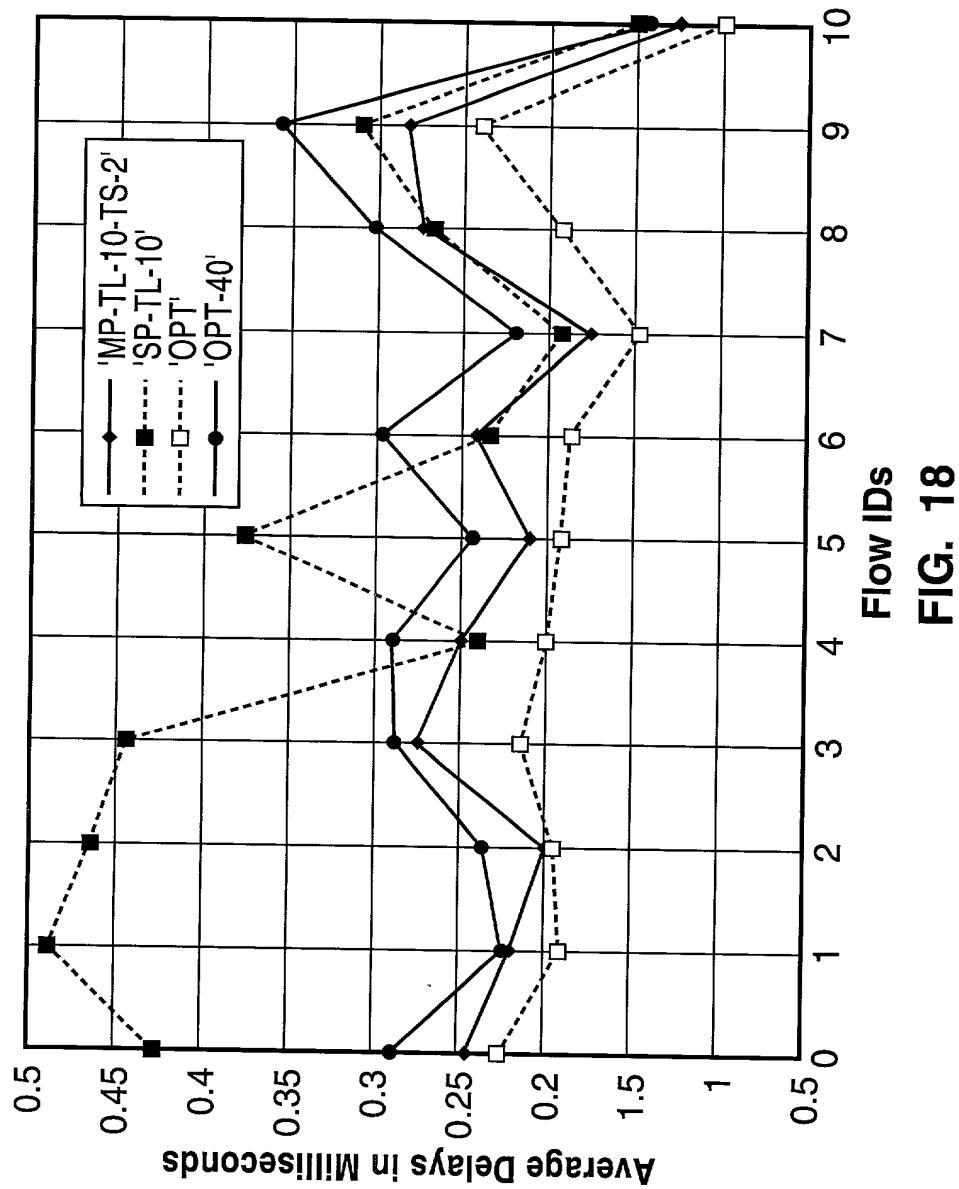


FIG. 18

FIG. 19